

Solution

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Problem 1:

Solution:

(a) No. The inequality is equivalent to

$$x_3 - \sqrt{(x_1 - y_1)^2 + (x_2 - y_2)^2 + (x_3 - y_3)^2} \leq 0.$$

However, the left-hand-side is a concave function. It is not hard to find a counter example.

(b) Yes.

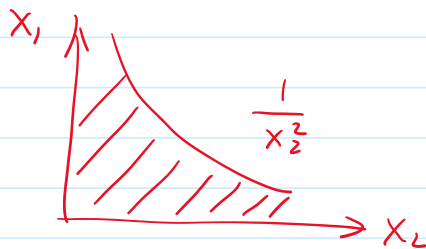
$$\underline{x} > a \Leftrightarrow v^T (\underline{x} - a \mathbb{1}) v \geq 0 \text{ for all } v$$

Thus, $\{\underline{x} \mid \underline{x} > a\}$

$$= \bigcap_v \{\underline{x} \mid v^T (\underline{x} - a \mathbb{1}) v \geq 0\}$$

The set $\{\underline{x} \mid v^T (\underline{x} - a \mathbb{1}) v \geq 0\}$ is convex since it is a half-space.

(c) No. It looks like this



which is clearly not convex.

(d) Yes.

p_1

(d) Yes.

$$\frac{p_1}{p_2 + N} > 2 \Leftrightarrow 2p_2 - p_1 + N < 0$$

It is a convex set since $2p_2 - p_1 + N$ is a convex (actually linear) function.

(e) Yes.

KL distance is a convex function in both \vec{p} & \vec{q} . This can be seen by the fact that $-p_i \log \frac{q_i}{p_i}$ is the perspective mapping of $-\log p_i$, which is convex.

Problem 2:

Solution: Let $z_1 \in A - B$. Then, there must exist $x_1 \in A, y_1 \in B$ such that $z_1 = x_1 - y_1$.

Similarly, let $z_2 \in A - B$, such that $z_2 = x_2 - y_2$, $x_2 \in A$ & $y_2 \in B$.

Consider the convex combination, $0 \leq \theta \leq 1$:

$$\begin{aligned} z &= z_1 \cdot \theta + z_2 (1 - \theta) \\ &= (x_1 - y_1) \theta + (x_2 - y_2) (1 - \theta) \\ &= [x_1 \theta + x_2 (1 - \theta)] - [y_1 \theta + y_2 (1 - \theta)] \end{aligned}$$

Since A is convex, $x_1 \theta + x_2 (1 - \theta) \in A$

Since B is convex, $y_1 \theta + y_2 (1 - \theta) \in B$

Thus, z must belong to $A - B$.

We can then conclude that $A - B$ is a convex set.

Problem 3:

(a) No. The first summation is equal to

$$\sum_{i: y_i=1} w^T x_i - \underbrace{\log(1 + e^{w^T x_i})}_{\text{convex}},$$

which is concave in w . Similarly, the second term is also concave

(b) No. When y is large,

$$\log\left(\frac{y^2}{x} + 1\right) \approx \log\frac{y^2}{x} = 2\log y - \log x$$

which is concave in y . Thus, the function cannot be convex in both x & y

(c) Yes. The function

$$\theta x - \log M_x(\theta)$$

is linear in x for every θ . Thus, $I(x)$ is convex in x since it is the pointwise maximum (over θ) of linear functions

(d) Yes. $\sqrt{xy} = x\sqrt{\frac{y}{x}}$ is concave since it is the

perspective mapping of \sqrt{y} .

Hence, $-\sqrt{x(x+y)}$ is a convex function.

(e) Yes. This follows from the composition rules.

- e^y is convex and increasing

- x^2 is convex

Problem 4:

Solution:

(a) The Lagrangian is

$$L(\vec{y}, \lambda) = \sum_{t=1}^T f(x_t + y_t) + \lambda \left(C - \sum_{t=1}^T y_t \right) \\ = \sum_{t=1}^T \left[f(x_t + y_t) - \lambda y_t \right] + \lambda C$$

Thus, the KKT condition for optimality is

(i) $\sum_{t=1}^T y_t \geq C, y_t \geq 0$

(ii) $\lambda \in \mathbb{R}$

(iii) y_t should minimize $f(x_t + y) - \lambda y$
over $y_t \geq 0$

(iv) $\lambda \left(\sum_{t=1}^T y_t - C \right) = 0$

(b) Consider any time-slot t such that $y_t > 0$.

The optimizing condition for (iii) states that

$$f'(x_t + y_t) - \lambda = 0$$

Since $f''(z) > 0$ for all z , the value of z such that $f'(z) = \lambda$ is unique.

Therefore, the value of $x_t + y_t$ must be the same for all time-slots such that $y_t > 0$.

Problem 5:

Solution:

(a) The Lagrangian is

$$L(a_1, a_2, \lambda_1, \lambda_2) = a_1 x_1 + a_2 x_2 + \lambda_1 (2a_1 + a_2 - 4) \\ + \lambda_2 (a_1 + 2a_2 - 5) \\ = a_1 (x_1 + 2\lambda_1 + \lambda_2) + a_2 (x_2 + \lambda_1 + 2\lambda_2) \\ - 4\lambda_1 - 5\lambda_2$$

To minimize L over a_1, a_2 , we get

To minimize L over a_1, a_2 , we get

$$\min_{a_1, a_2} L(a_1, a_2, \lambda_1, \lambda_2) = \begin{cases} -4\lambda_1 - 5\lambda_2, & \text{if } \begin{cases} x_1 + 2\lambda_1 + \lambda_2 = 0 \\ x_2 + \lambda_1 + 2\lambda_2 = 0 \end{cases} \\ -\infty, & \text{o/w.} \end{cases}$$

Therefore, the dual problem is

$$\begin{aligned} \max_{\lambda_1, \lambda_2} & -4\lambda_1 - 5\lambda_2 && (7) \\ \text{sub to} & x_1 + 2\lambda_1 + \lambda_2 = 0 && (8) \\ & x_2 + \lambda_1 + 2\lambda_2 = 0 \\ & \lambda_1 \geq 0, \lambda_2 \geq 0 \end{aligned}$$

(b) By strong duality.

$$\min_{a \in P} a^T x = \max_{\lambda_1, \lambda_2 \in (8)} -4\lambda_1 - 5\lambda_2$$

Thus, the constraint that

$$\min_{a \in P} a^T x \geq 1$$

$$a \in P$$

is equivalent to

$$\max_{\lambda_1, \lambda_2 \in (8)} -4\lambda_1 - 5\lambda_2 \geq 1$$

$$\lambda_1, \lambda_2 \in (8)$$

and is equivalent to

There exists $\lambda_1, \lambda_2 \in (8)$ such

$$\text{that } -4\lambda_1 - 5\lambda_2 \geq 1$$

Therefore, (1) is equivalent to

$$\begin{aligned} \min & f_0(x) \\ \text{sub to} & -4\lambda_1 - 5\lambda_2 \geq 1 \\ & x_1 + 2\lambda_1 + \lambda_2 = 0 \\ & x_2 + \lambda_1 + 2\lambda_2 = 0 \\ & \lambda_1 \geq 0, \lambda_2 \geq 0 \end{aligned}$$

Problem 6:

Solution:

(a) The only non-convex constraint is (10).
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(a) The only non-convex constraint is (10).

However, we can replace it by

$$r_l \leq \log\left(1 + \frac{\partial L P_l}{w_0}\right) \quad (11)$$

because the optimal solution will still pick the largest r_l , which will make (11) tight with equality.

- Associate the Lagrange multiplier λ_l .

The Lagrangian is

$$\begin{aligned} L(\vec{x}, \vec{r}, \vec{p}, \vec{\lambda}) &= -\sum_s U_s(x_s) + \sum_l P_l \\ &\quad + \sum_l \lambda_l \left[\sum_s H_s^l x_s - r_l \right] \\ &= -\sum_s \left[U_s(x_s) - x_s \cdot \sum_l H_s^l \lambda_l \right] \\ &\quad + \sum_l \left[P_l - \lambda_l r_l \right] \end{aligned}$$

- Therefore, to minimize L over $\vec{x}, \vec{r}, \vec{p}$ given $\vec{\lambda}$, each flow s can solve

$$x_s(t) = \max U_s(x_s) - x_s \sum_l H_s^l \lambda_l(t) \quad (12)$$

Each link can solve

$$P_l(t) = \min P_l - \lambda_l r_l \quad (13)$$

$$\text{sub to } r_l \leq \log\left(1 + \frac{\partial L P_l}{w_0}\right)$$

$$\Delta r_l(t) = \log\left(1 + \frac{\partial L P_l(t)}{w_0}\right).$$

- To solve (12), we can set the derivative to zero (assuming that $x_s \geq 0$):

$$U_s'(x_s) = \sum_l H_s^l \lambda_l(t)$$

$$\Rightarrow x_s(t) = \max\{0, U_s'^{-1}\left(\sum_l H_s^l \lambda_l(t)\right)\}$$

To solve (13), we note that it is equivalent to

To solve (13), we note that it is equivalent to

$$\min_{P_c} P_c - \lambda_c \log \left(1 + \frac{\partial_c P_c}{W_0} \right)$$

Assuming $P_c > 0$, we can also set the derivative to zero:

$$1 - \lambda_c \frac{\frac{\partial_c}{W_0}}{1 + \frac{\partial_c P_c}{W_0}} = 0$$

$$\Rightarrow P_c(t) = \max \left\{ 0, \lambda_c - \frac{W_0}{\partial_c} \right\}.$$

- Finally, the dual algorithm can update

$$\lambda_c \text{ by } \lambda_c(t+1) = \left\{ \lambda_c(t) + \gamma \left[\sum_S H_S^L X_S(t) - r_c(t) \right] \right\}^+ \quad (14)$$

(b) Indeed, the above dual algorithm indeed represents a distributed & price-driven operation: Given λ_c

- each flow can solve X_S independently based on (12)
- each link can determine P_c independently based on (13).
- each link can update λ_c based on (14), without knowledge of $U_S(\cdot)$ & rate-power function.

Note: For problem (c), some students update x, p, λ all in the gradient direction:

$$x_S(t+1) = x_S(t) - \alpha \frac{\partial L}{\partial x_S}$$

$$p_i(t+1) = p_i(t) - \alpha \frac{\partial L}{\partial p_i}$$

$$\lambda_i(t+1) = \lambda_i(t) + \alpha \frac{\partial L}{\partial \lambda_i}$$

However, note that such updates neither maximize nor minimize $L(x, p, \lambda)$. Thus, the convergence and optimality of this type of "primal-dual" algorithm is not obvious.

Instead, this problem asks for the "dual" algorithm, which is stated above. It is optimal since it maximizes the dual objective function.